

# Towards LLM-support for Deductive Verification of Java Programs

Theorem Proving and Machine Learning in the Age of LLMs

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**Samuel Teuber**, Bernhard Beckert

Institute of Information Security and Dependability (KASTEL)

2025

# Motivation: Can LLMs Generate Specifications?

Large Language Models have seen tremendous success in recent years

GitHub Copilot & Co show: LLMs can **generate code**

But can they **specify code**?

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  @*/
public static /*@ pure */ int max(int[] a) {
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  int max = a[0], i = 1;
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  while (i < a.length) {
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    ++i;
  }
  return max;
}
```

## Verification requires Loop Invariant

- Holds before first loop iteration
- Preserved by loop iteration
- Implies post condition

## Additionally:

- Loop *Variant*
- Assignable Heap Variables

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    @ loop_invariant (\exists int k; 0 <= k && k < i; max == a[k]);
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Writing auxiliary spec yourself [ChatGPT]

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Having ChatGPT write auxiliary spec [ChatGPT]

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## Let's ask ChatGPT:

Do you know JML, the Java Modeling Language?

Yes, I am familiar with JML (Java Modeling Language). JML is a formal specification language for Java programs. [...] JML is typically used in conjunction with formal verification tools, such as ESC/Java or **KeY**, to check that the code meets its specifications.

# The Program Verifier KeY

Deductive verification

100% Java Card

Java Modeling  
Language (JML)



## Numerous Case Studies:

- TimSort (OpenJDK)
- `LinkedList` (OpenJDK)
- Super Scalar Sample Sort

Modular  
Reasoning

Verification Methodology:

Theorem Proving for **Java Dynamic Logic**

collaboration of KIT, TU Darmstadt, Chalmers University

Ahrendt u. a. 2016



# Java Dynamic Logic by Example

```
class MyClass {
  int arr [];
  int maxOfFirstTwo() {
    if (this.arr[0] > this.arr[1]) {
      return this.arr[0];
    } else {
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JavaDL allows reasoning about **real-world Java** programs.

We can verify programs by **proving the validity of formulas via a calculus**:

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```
res = maxOfFirstTwo();
```

Java Program

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## Proof Interaction:

- Interactive (manual) application of proof rules

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## Proof Interaction:

- Interactive (manual) application of proof rules
- Annotation of source code **auto-active verification** (Leino und Moskal 2010)

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# Java Modelling Language

- Specification Language for Java
- Design by Contract Paradigm
- Rich set of possible first-order annotations:
  - Hoare-Style pre- and post-conditions
  - Invariants
  - Asserts
  - Class-Invariants
- Supported by numerous tools for Java verification

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# LLMs for Deductive Java Verification

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- May produce output that **is not correct**
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## Objective: An Intersymbolic AI approach to Program Verification

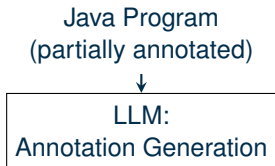
Combine LLMs and Deductive Verification so that **weaknesses cancel out**

# LLM-based generation of JML: Approach

Java Program  
(partially annotated)

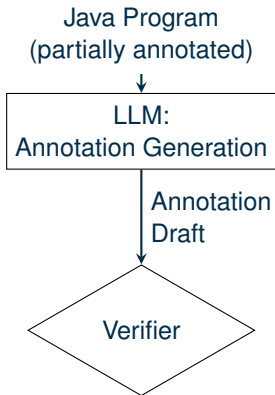
```
//@ ensures \result == -2*x;  
int f(int x) {  
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}  
  
int g(int x) {  
    return x+x;  
}
```

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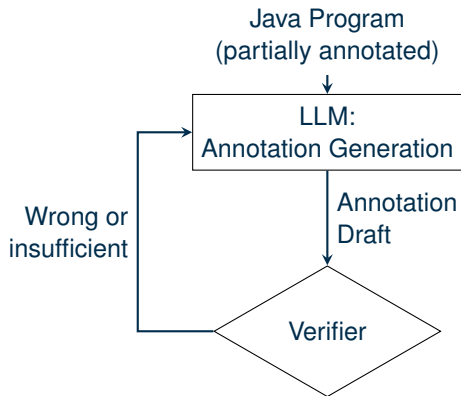
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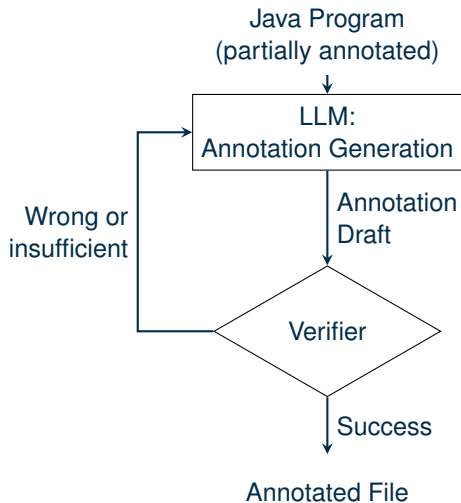
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/*@ ensures x == 2 ==> \result == 4;  
int g(int x) {  
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# Evaluation

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KeY repository and old exercise sheets



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- Generate auxiliary annotations for given top-level spec  
(Java+JML  $\Rightarrow$  JML)
  - Loop invariant (given method contract)
  - Contract of callee-method (given caller contract)

# Evaluation Benchmark: Features

	Isolated Methods	Sub-Methods	Invariants	Total	
<b>No. of benchmarks</b>	36	27	14	77	
<b>JML features</b>					
Quantifiers	32	24	14	70	91%
Non-empty assignable	24	17	11	52	68%
Reference to pre-state (\old)	18	11	8	37	48%
(Pure) Method calls in spec	10	7	7	24	31%
<b>Java features</b>					
Arrays	30	23	13	66	86%
Field access	11	7	3	21	27%

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Category	# Benchmarks	$\mu \pm \sigma$ of success rate (%)	
		GPT 3.5 ( $n = 10$ )	GPT 4o ( $n = 3$ )
Isolated Method	36	52.2 $\pm$ 4.3	<b>62.0 <math>\pm</math> 1.6</b>
Submethods	14	19.3 $\pm$ 12.1	<b>40.5 <math>\pm</math> 4.1</b>
Invariants	27	37.0 $\pm$ 7.4	<b>67.9 <math>\pm</math> 5.7</b>

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## No elaborate prompt engineering yet!

- **Objective:** Evaluate baseline performance
- Not enough benchmarks

# Evaluation: Isolated Methods

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**Success criterion: KeY proves code satisfies generated spec**

**Manual Inspection: Spec adequately specifies code behavior**  
Sometimes incomplete

# Evaluation: Isolated Methods – Errors

Error Category	Share of error (%)	
	GPT 3.5	GPT 4o
<b>Syntactic Errors</b>		
Loop Invariant Generation	4.3	1.3
Unknown variable names	11.8	2.2
Incorrect usage of <code>\result</code>	1.1	0.3
Other parsing errors	17.1	10.7
<b>Semantic Errors</b>		
Incomplete proof	49.5	70.3
Timeout of Verifier	16.0	15.1

## Observation

GPT 4o tends to make **semantic** errors (much less syntactic errors).

# Demo

Backup



# Grasping Feedback

We can provide feedback on failed proof attempts,  
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The provided invariant **does not hold for the following instantiations of variables:**

k = 0  
a = [0, 0, 0]  
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You're right; we still need to refine the loop invariant to handle the specific case where **the array length is odd**. In such cases, the middle element should remain unchanged.

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Here's the corrected code:  
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**Sampling (pass@N) vs. Feedback**  
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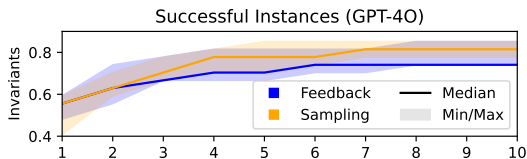
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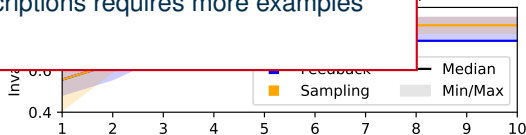
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**Focus:** Auxilliary Specifications

**We need more benchmarks for conclusive results**

Prompt Engineering for proof state descriptions requires more examples



# Feedback vs. Sampling

## What is the right metric?

### Classic Verification

Two verification techniques:

(A) 5 iterations, 2 seconds CPU time/iteration

(B) 3 iterations, 4 seconds CPU time/iteration

⇒ **(A) is better**



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### What we know:

Computational cost increases with token count

Initial Query:  $I$  tokens

LLM Output:  $O$  tokens

Feedback:  $F$  tokens

**Sampling:**  $n(I + O) \in \mathcal{O}(n)$  tokens

**Feedback:**  $n(I + O) + \frac{n(n-1)}{2}(O + F) \in \mathcal{O}(n^2)$  tokens

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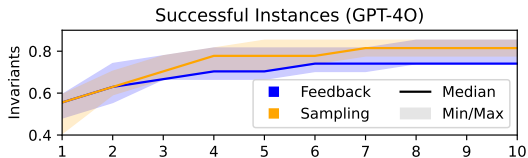
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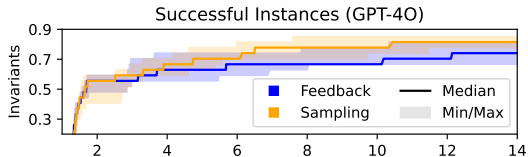
### Another evaluation:

Iterations → Normalized Token Count ( $\frac{\# \text{Tokens}}{I}$ )

### Iteration based:



### Token based:



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Strategy for generating all method/invariant annotations
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Loughridge u. a. [2025](#)

Wen u. a. [2024](#)

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Dataset for JML via GitHub scraping

Greiner u. a. [2024](#)



# Conclusion

- OpenAI's models are **surprisingly good** at JML
- Iteration helps, but feedback not (yet)
- **Foundation Models:**
  - Make effective usage of ML for niche languages like JML possible
  - Cannot solve all data problems (we still need well-curated benchmark sets)

## Open Questions

- Effective Feedback from the theorem prover
- Prompt Engineering in the presence of “*perfect checkers*”  
Sampling may not be that bad afterall?
- What is a fair comparison between an LLM approach and a “classic” approach?

# References I

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# Demo (Backup)

Back to Slides

# Example: Find the Maximum

```
/*@ normal_behavior
   @ ensures (\forall int j; j >= 0 && j < a.length; \result >= a[j]);
   @ ensures a.length > 0 ==>
   @ (\exists int j; j >= 0 && j < a.length; \result == a[j]);
   @ */
public static /*@ pure */ int max(int[] a) {
    if (a.length == 0) return 0;
    int max = a[0], i = 1;
    while (i < a.length) {
        if (a[i] > max) max = a[i];
        ++i;
    }
    return max;
}
```

# Example: Find the Maximum (1st try)

Given the following Java class:

[ ... ]

Please **provide a loop invariant** for the loop construct with the **comment** **'//Add invariant here'** of the method 'max'. Beware this annotation has to be a loop invariant. A loop invariant **typically has the following structure**:

[Template: **loop\_invariant, decreases, assignable** ]

Based on the method 'max', the goal is to find the maximum element in the given array 'a'. [...]

```
/*@ loop_invariant 0 <= i && i <= a.length;  
   @ loop_invariant (\forall int k; 0 <= k && k < i; max  
       >= a[k]);  
   @ decreases a.length - i;  
   @ assignable i, max;  
   @ */
```

**Problem:** max not required to be in the array

## Example: Find the Maximum (2nd try)

The provided JML does not solve the task. [...] **some proof goals were not closed:**

[...]

During verification, the following proof branches could not be closed:

**Normal Execution ( $\_a \neq \text{null}$ )**

Please fix the JML loop invariant.

Let's refine the loop invariant [...]

```
/*@ loop_invariant a != null;
   @ loop_invariant 0 <= i && i <= a.length;
   @ loop_invariant (\forall int k; 0 <= k && k < i; max
                     >= a[k]);
   @ decreases a.length - i;
   @ assignable i, max;
  @*/
```

**Problem:** same (only added `a != null`)

## Example: Find the Maximum (3rd try)

[Same message]

To refine the loop invariant further [...]

```
/*@ loop_invariant a != null;  
  @ loop_invariant 0 <= i && i <= a.length;  
  @ loop_invariant (\forall int k; 0 <= k && k < i; max  
    >= a[k]);  
  @ loop_invariant (i > 0) ==>  
  @ (\exists int k; 0 <= k && k < i; max == a[k]);  
  @ decreases a.length - i;  
  @ assignable i, max;  
  @*/
```

Invariant that is strong enough to prove the post-condition.

**Problem:** Invariant is not inductive (because of  $i > 0$ )



# Example: Find the Maximum (4th try)

[Same message]

Let's ensure that the JML loop invariant comprehensively captures [...]

```
/*@ loop_invariant a != null;
   @ loop_invariant 0 <= i && i <= a.length;
   @ loop_invariant (\forall int k; 0 <= k && k < i; max
   >= a[k]);
   @ loop_invariant
   @ \exists int m; m >= 0 && m < i && max == a[m];
   @ decreases a.length - i;
   @ assignable i, max;
   @ */
```

Correct invariant

**Syntax error:** ACSL `\exists ...` vs. JML (`\exists ...`)

## Example: Find the Maximum (5th try)

The provided code is not valid JML.

[...]

Error during JML parsing: Failed to parse JML fragment: **Encountered unexpected token: “\exists” “\exists” at line 4, column 20.**

Was expecting one of: [...]

[unfortunately still fails]

# Example: Find the Maximum (restart)

## Second Attempt...

[Same initial question]

In the given 'max' method, the loop is responsible for [...]

```
/*@ loop_invariant 0 <= i && i <= a.length;  
   @ loop_invariant (\forall int k; 0 <= k && k < i; max  
     >= a[k]);  
   @ decreases a.length - i;  
   @ assignable max, i;  
  @*/
```

**Problem:** `max` not required to be in array.

[Same feedback]

[Correct answer]

# More Difficult Example: Rotate an Array

```
public static void rotate(int[] a, int len) {
    int[] b = new int[a.length];
    int i = 0;
    /*@ [...] @*/
    for (i = 0; i < len; i++) {
        b[i] = a[a.length - len + i];
    }
    /*@ [...] @*/
    for (i = len; i < a.length; i++) {
        b[i] = a[i - len];
    }
    /*@ [...] @*/
    for (i = 0; i < a.length; i++) {
        a[i] = b[i];
    }
}
```

# More Difficult Example: Rotate an Array

## Correct specification generated by GPT 4o

```
/* @ normal_behavior
   @
   @ requires a != null && 0 <= len && len <= a.length;
   @
   @ assignable a[*];
   @
   @ ensures (\forall int i; 0 <= i && i < len;
             @           a[i] == \old(a[a.length - len + i]));
   @ ensures (\forall int i; len <= i && i < a.length;
             @           a[i] == \old(a[i - len]));
   @*/
```

# Demo (Backup)

Back to Slides

# Evaluation: Isolated Methods

Category	# Benchmarks	$\mu \pm \sigma$ of success rate (%)	
		GPT 3.5 ( $n = 10$ )	GPT 4o ( $n = 3$ )
Isolated Method	36	52.2 $\pm$ 4.3	<b>62.0 <math>\pm</math> 1.6</b>

**Success criterion: KeY proves code satisfies generated spec**

**Manual Inspection: Spec adequately specifies code behavior**  
Sometimes incomplete

## Repetition/Feedback helps

- 75% of benchmarks successful (over 10 runs, GPT 3.5)
- Feedback from the verifier can help

